

CLASSIFICATION OF TWO-REGULAR DIGRAPHS WITH MAXIMUM DIAMETER

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ABSTRACT. The Klee-Quaife problem is finding the minimum order $\mu(d, c, v)$ of the (d, c, v) graph, which is a c -vertex connected v -regular graph with diameter d . Many authors contributed finding $\mu(d, c, v)$ and they also enumerated and classified the graphs in several cases. This problem is naturally extended to the case of digraphs. So we are interested in the extended Klee-Quaife problem. In this paper, we deal with an equivalent problem, finding the maximum diameter of digraphs with given order, focused on 2-regular case. We show that the maximum diameter of strongly connected 2-regular digraphs with order n is $n - 3$, and classify the digraphs which have diameter $n - 3$. All 15 nonisomorphic extremal digraphs are listed.

1. Introduction

Let G be a connected graph. G is c -vertex connected if the graph obtained by deleting arbitrary $c - 1$ vertices from G remains connected. G is v -regular if each vertex of G is adjacent to exactly v vertices. A (d, c, v) graph (resp. $\langle d, c, v \rangle$ graph) is a c -vertex connected v -regular (resp. minimum degree v) graph with diameter d . The number $\mu(d, c, v)$ (resp. $\mu\langle d, c, v \rangle$) is the minimum order of the (d, c, v) (resp. $\langle d, c, v \rangle$) graphs and a minimum (d, c, v) graph is a (d, c, v) graph on $\mu(d, c, v)$ vertices.

In 1960's, there have been some early results which are equivalent to computing $\mu(d, c, v)$ and $\mu\langle d, c, v \rangle$ for special cases [4, 6, 10]. In 1976,

Received April 29, 2012. Revised June 13, 2012. Accepted June 15, 2012.

2010 Mathematics Subject Classification: 05C20, 05C12, 05C07.

Key words and phrases: 2-regular, diameter, digraphs.

This work was supported by the Research Institute of Natural Science of Gangneung-Wonju National University.

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concerning with armed connection network, Klee and Quaife [8] raised a problem of finding $\mu(d, c, v)$ and $\mu\langle d, c, v \rangle$. In the same paper, $\mu\langle d, c, v \rangle$ was computed completely. For $\mu(d, c, v)$, there are some partial results and it was lastly computed in 1989 for all d, c, v [3].

Classifying and enumerating all minimum (d, c, v) graphs is a more complicated problem. This problem was solved only in two cases. The first case is $v = 3$ and $c = 1, 2$ done by Klee and Quaife [9] in 1977. The second case is $v = c = 3$ and d is odd. This was achieved by Klee [7] in 1980. It is also notable Mayers [11] found a method to construct all the $(d, 3, 3)$ graphs for all d in 1980, whereas he couldn't enumerate all $(d, 3, 3)$ graphs. Bhattacharya found a method to construct some minimum (d, n, n) graphs in 1985 [1]. But his method didn't cover all minimum (d, n, n) graphs.

Now we consider the corresponding problem for digraph $D = (V, A)$. A digraph D is strongly connected if for each pair of vertices x, y of D there is a directed walk from x to y . The ingegree $\delta^+(x)$ (respectively, the outdegree $\delta^-(x)$) of a vertex x in D is the number of vertices y in D such that $(y, x) \in A$ (respectively, $(x, y) \in A$). A digraph D is eulerian if $\delta^+(x) = \delta^-(x)$ for each vertex x in D . D is oriented if there is no pair of vertices x, y in D such that $(x, y) \in A$ and $(y, x) \in A$. The (d, c, δ) digraphs and $\langle d, c, \delta \rangle$ eulerian digraphs are defined similarly as the case of graphs. Finding minimum order of (d, c, δ) digraphs and $\langle d, c, \delta \rangle$ eulerian digraphs is equivalent to determining maximum diameter d such that there is (d, c, v) and $\langle d, c, v \rangle$ eulerian digraphs on μ vertices, respectively. In [5], Knyazev proved that if D is an eulerian oriented digraph on n vertices, then $\frac{4n}{2\delta+1} - 4 \leq \text{diam}(D) \leq \frac{5}{2\delta+n}n$. Dankelmann [2] improved the upper bound of $\text{diam}(D)$ to $\frac{4}{2\delta+1}n + 2$. Their results imply that the minimum order μ of an eulerian oriented $\langle d, 1, \delta \rangle$ digraph satisfies $\frac{(2\delta+1)(d-2)}{4} \leq \mu \leq \frac{(2\delta+1)(d+4)}{4}$.

In this paper, we show that the maximum diameter of strongly connected 2-regular digraphs on n vertices is $n - 3$ when $n \geq 9$. As a consequence we have the maximum number of vertices in a strongly connected oriented eulerian $(d, 1, 2)$ digraph is $d + 3$ when $d \geq 6$. Moreover, in this case we classify all 15 digraphs of diameter $n - 3$.

2. Main theorems

Let $D = (V, A)$ be a strongly connected digraph on n vertices. Assume that D is 2-regular. For a vertex v in D , we define $A^+(v) = \{w \in V \mid (v, w) \in A\}$ and $A^-(v) = \{w \in V \mid (w, v) \in A\}$. We also define $v \xrightarrow{k} w$ for each pair of vertices v, w in D if there is a walk of length k from v to w . We use $v \longrightarrow w$ instead of $v \xrightarrow{1} w$. Since D is 2-regular, the outdegree $\delta^+(v)$ and indegree $\delta^-(v)$ are 2 for every vertex $v \in V$.

THEOREM 1. *If D is a strongly connected 2-regular digraph on $n(\geq 5)$ vertices, then $\text{diam}(D) \leq n - 3$.*

Proof. Suppose that $\text{diam}(D) \geq n - 2$. There are vertices v, w in D such that $\text{dist}(v, w) = n - 2$. Since $v \xrightarrow{n-2} w$, there are v_0, v_1, \dots, v_{n-2} such that $v = v_0 \rightarrow v_1 \rightarrow v_2 \rightarrow \dots \rightarrow v_{n-2} = w$. Since $\delta^+(v) = 2$ and $\delta^-(w) = 2$, there are vertices $v'_1 (\neq v_1), v'_{n-3} (\neq v_{n-3})$ such that $v'_1 \in A^+(v)$ and $v'_{n-3} \in A^-(w)$. Since $v'_1 \neq v_0, v_1$ and $\text{dist}(v, w) = n - 2$, $v'_1 \notin \{v_0, v_1, \dots, v_{n-2}\}$. Similarly, we can show that $v'_{n-3} \notin \{v_0, v_1, \dots, v_{n-2}\}$. If $v'_1 = v'_{n-3}$, then $v \xrightarrow{2} w$ and $\text{dist}(v, w) = 2 < n - 2$, which is a contradiction. So the vertex set V includes $\{v_0, v_1, \dots, v_{n-2}, v'_1, v'_{n-3}\}$, which contradicts $|V| = n$. So $\text{diam}(D) \leq n - 3$. □

By the above theorem, $\text{diam}(D) \leq n - 3$. From now on we assume that $\text{diam}(D) = n - 3$. Then there are vertices v, w such that $\text{dist}(v, w) = n - 3$. So there are vertices v_0, v_1, \dots, v_{n-3} such that $v = v_0 \rightarrow v_1 \rightarrow v_2 \rightarrow \dots \rightarrow v_{n-3} = w$. And there are vertices $v'_1 (\neq v_1), v'_{n-4} (\neq v_{n-4})$ such that $v_0 \longrightarrow v'_1$ and $v'_{n-4} \longrightarrow v_{n-3}$. In this case, $V = \{v_0, v_1, \dots, v_{n-3}, v'_1, v'_{n-4}\}$. Using these notations, we have the following Lemma.

LEMMA 1. *Let $n \geq 7$ and $1 \leq i \leq n - 6$. If $x \in \{v_0, v_1, \dots, v_i, v'_1\}$ and $y \in \{v_{i+2}, \dots, v_{n-3}, v'_{n-4}\}$, then $(x, y) \notin A$.*

Proof. If $(x, y) \in A$, then

$$\begin{aligned} n - 3 = \text{dist}(v_0, v_{n-3}) &\leq \text{dist}(v_0, x) + \text{dist}(x, y) + \text{dist}(y, v_{n-3}) \\ &\leq i + 1 + n - i - 5 = n - 4. \end{aligned}$$

This is a contradiction. □

COROLLARY 1. *If $n \geq 10$ and $3 \leq i \leq n - 7$, then $v_{i+1} \longrightarrow v_i$.*

Proof. Let $A_i = \{v_0, v_1, \dots, v_i, v'_1\}$ and $B_i = \{v_{i+1}, \dots, v_{n-3}, v'_{n-4}\}$. Since $A_i \cup B_i = V$, $A_i \cap B_i = \phi$ and D is 2-regular, the number of arcs from A_i to B_i and from B_i to A_i are equal. By Lemma 1, there are no arc from A_{i-1} to B_i and from A_i to B_{i+1} . So (v_i, v_{i+1}) is the only arc from A_i to B_i . Thus there is only one arc from B_i to A_i . Let $A^+(v_{i+1}) = \{v_{i+2}, x\}$ and $A^-(v_i) = \{v_{i-1}, y\}$. Since there is no arc from $\{v_{i+1}\}$ to B_{i+2} , $x \in A_i$. Since (v_{i+1}, x) is the only arc from B_i to A_i , there is no arc from B_{i+1} to $\{v_i\}$. So $y \in A_{i+1}$. By Lemma 1, there is no arc from A_{i-2} to $\{v_i\}$, which implies $y \notin A_{i-2}$. Since $y \neq v_{i-1}, v_i$ and $y = v_{i+1}$. Thus, $v_{i+1} \rightarrow v_i$. \square

By Corollary 1, $v_{i+1} \rightarrow v_i \rightarrow v_{i-1}$ for $n \geq 11$ and $4 \leq i \leq n-7$. Since D is 2-regular we have the following corollary.

COROLLARY 2. *If $n \geq 11$ and $4 \leq i \leq n-7$, then $A^-(v_i) = A^+(v_i) = \{v_{i-1}, v_{i+1}\}$.*

If $V' \subset V$, then we use $\langle V' \rangle$ to be the directed subgraph of $D = (V, A)$ which is induced by V' . For $n \geq 2$, let P_n be the path on n vertices and P_1 be trivial graph.

LEMMA 2. *If $n \geq 9$, then $\langle v_3, \dots, v_{n-6} \rangle$ is isomorphic to the path P_{n-8} .*

Proof. It is trivial when $n = 9, 10$. When $n \geq 11$, by Lemma 1, $v_3 \rightarrow v_{n-6}$. Since $A^+(v_{n-6}) = \{v_{n-5}, v_{n-7}\}$, $v_{n-6} \rightarrow v_3$. So by Corollary 2, $\langle v_3, \dots, v_{n-6} \rangle$ is isomorphic to the path P_{n-8} . \square

LEMMA 3. *If $n \geq 9$, then we have the followings.*

- (1) $\langle v_0, \dots, v_3, v'_1 \rangle$ is isomorphic to one of H_1, \dots, H_5 in Figure 1.
- (2) $\langle v_{n-6}, \dots, v_{n-3}, v'_{n-4} \rangle$ is isomorphic to one of T_1, \dots, T_5 in Figure 1.

Proof. We divide the proof into the cases according to how the sets $A^+(v_1), A^+(v'_1)$ are given.

Case1. $A^+(v_1) = \{v_0, v_2\}, A^+(v'_1) = \{v_0, v_1\}$

Since $A^-(v_0) = \{v'_1, v_1\}$ and $A^-(v_1) = \{v_0, v'_1\}$, $v_2 \rightarrow v'_1$. Since $A^-(v'_1) = \{v_0, v_2\}$, by Lemma 1, $v_3 \rightarrow v_2$. So $\langle v_0, v_1, v_2, v_3, v'_1 \rangle$ is isomorphic to H_1 in Figure 1.

Case2. $A^+(v_1) = \{v'_1, v_2\}, A^+(v'_1) = \{v_0, v_1\}$

Since $A^-(v'_1) = \{v_0, v_1\}$ and $A^-(v_1) = \{v_0, v'_1\}$, by Lemma 1 $v_2 \rightarrow v_0$. Lemma 1 and the fact that $A^-(v_0) = \{v'_1, v_2\}$ imply $v_3 \rightarrow v_2$. We can

see $\langle v'_1, v_1, v_2, v_3, v_0 \rangle$ is isomorphic to H_1 in Figure 1.

Case3. $A^+(v_1) = A^+(v'_1) = \{v_0, v_2\}$

Since $A^-(v_0) = \{v_1, v'_1\}$, $v_2 \rightarrow v_0$. By Lemma 1, $v_2 \rightarrow v_1$ or $v_2 \rightarrow v'_1$. If $v_2 \rightarrow v_1$, $A^-(v_1) = \{v_0, v_2\}$. Since $A^-(v_2) = A^-(v_0) = \{v_1, v'_1\}$, by Lemma 1 $v_3 \rightarrow v'_1$. So $\langle v_0, v_1, v_2, v_3, v'_1 \rangle = H_2$ in Figure 1. If $v_2 \rightarrow v'_1$, then we have $v_3 \rightarrow v_1$. And $\langle v_0, v'_1, v_2, v_3, v_1 \rangle$ is isomorphic to H_2 in Figure 1.

Case4. $A^+(v_1) = \{v'_1, v_2\}$, $A^+(v'_1) = \{v_0, v_2\}$

Since $A^-(v'_1) = \{v_0, v_1\}$, $v_2 \rightarrow v'_1$. So $v_2 \rightarrow v_1$ or $v_2 \rightarrow v_0$. If $v_2 \rightarrow v_0$, since $A^-(v_0) = \{v'_1, v_2\}$ and $A^-(v_2) = \{v'_1, v_1\}$, by Lemma 1 $v_3 \rightarrow v_1$. So $\langle v_0, v_1, v_2, v_3, v'_1 \rangle$ is isomorphic to H_3 in Figure 1. If $v_2 \rightarrow v_1$, then $A^-(v'_1) = \{v_0, v_1\}$, $A^-(v_1) = \{v_0, v_2\}$ and $A^-(v_2) = \{v_1, v'_1\}$. By Lemma 1, $v_3 \rightarrow v_0$. So $\langle v_0, v_1, v_2, v_3, v'_1 \rangle$ is isomorphic to H_4 in Figure 1.

Case5. $A^+(v_1) = \{v_0, v_2\}$, $A^+(v'_1) = \{v_1, v_2\}$

Since $A^-(v_1) = \{v_0, v'_1\}$, by Lemma 1 $v_2 \rightarrow v_0$ or $v_2 \rightarrow v'_1$. If $v_2 \rightarrow v_0$, since $A^-(v_0) = \{v_1, v_2\}$, $A^-(v_1) = \{v_0, v'_1\}$ and $A^-(v_2) = \{v_1, v'_1\}$, by Lemma 1 $v_3 \rightarrow v'_1$. So $\langle v_0, v'_1, v_2, v_3, v_1 \rangle$ is isomorphic to H_3 in Figure 1. If $v_2 \rightarrow v'_1$, we must have $v_3 \rightarrow v_0$. We can see $\langle v_0, v'_1, v_2, v_3, v_1 \rangle$ is isomorphic to H_4 in Figure 1. Case6. $A^+(v_1) = \{v'_1, v_2\}$, $A^+(v'_1) = \{v_1, v_2\}$

Lemma 1, $A^-(v_1) = \{v_0, v'_1\}$, $A^-(v'_1) = \{v_0, v_1\}$ and $A^-(v_2) = \{v_1, v'_1\}$ imply $v_2 \rightarrow v_0$ and $v_3 \rightarrow v_0$. So $\langle v_0, v_1, v_2, v_3, v'_1 \rangle$ is isomorphic to H_5 in Figure 1. So (1) holds. Similarly we can prove (2) by substituting v_0, v_1, v_2, v_3, v'_1 with $v_{n-3}, v_{n-4}, v_{n-5}, v_{n-6}, v'_{n-4}$ respectively. □

We call H_1, \dots, H_5 in Lemma 3 as heads and T_1, \dots, T_5 in Lemma 3 as tails. We can see that the union of $\langle v_0, v_1, v_2, v_3, v'_1 \rangle$, $\langle v_3, \dots, v_{n-6} \rangle$ and $\langle v_{n-6}, v_{n-5}, v_{n-4}, v_{n-3}, v'_{n-4} \rangle$ is a 2-regular digraph on V . So D is the union of subgraphs $\langle v_0, v_1, v_2, v_3, v'_1 \rangle$, $\langle v_3, \dots, v_{n-6} \rangle$ and $\langle v_{n-6}, v_{n-5}, v_{n-4}, v_{n-3}, v'_{n-4} \rangle$. Let $D_{i,j}$ be the union $H_i \cup P_{n-8} \cup T_j$ of digraphs H_i, P_{n-8} and T_j for $1 \leq i, j \leq 5$.

THEOREM 2. *If D is a strongly connected 2-regular digraph on n vertices and $\text{diam}(D) = n - 3$, then D is isomorphic to one of $\{D_{i,j} | 1 \leq i \leq j \leq 5\}$ in Figure 2.*

Proof. Let $D_{i,j} = H_i \cup P_{n-8} \cup T_j$. The functions f_i defined by

$$(f_i(v_0), f_i(v_1), f_i(v_2), f_i(v_3), f_i(v'_1)) = F_i$$

where $F_1 = (v_{n-3}, v'_{n-4}, v_{n-5}, v_{n-6}, v_{n-4})$, $F_2 = (v_{n-4}, v_{n-3}, v'_{n-4}, v_{n-6}, v_{n-5})$, $F_3 = (v_{n-3}, v_{n-5}, v_{n-4}, v_{n-6}, v'_{n-4})$, $F_4 = (v_{n-5}, v'_{n-4}, v_{n-4}, v_{n-6}, v_{n-4})$,

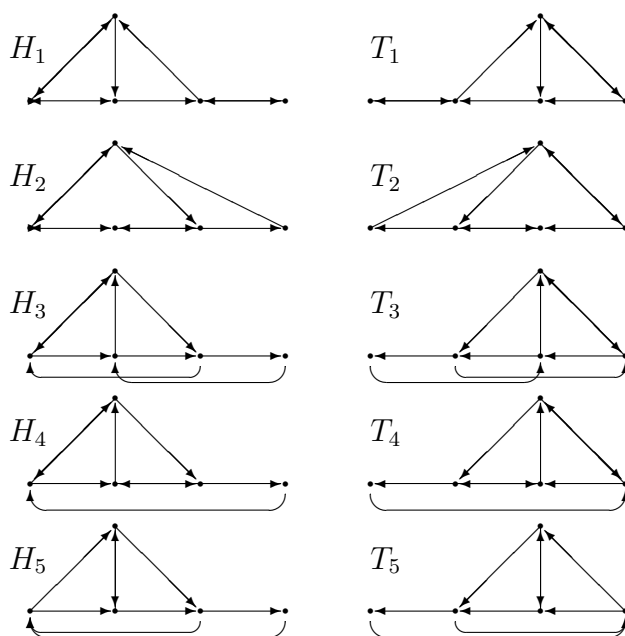


Figure 1. Heads and tails

and $F_5 = (v_{n-5}, v_{n-4}, v_{n-3}, v_{n-6}, v'_{n-4})$ give isomorphisms from H_i to T_i for all $i = 1, 2, \dots, 5$. Thus $D_{i,j}$ and $D_{j,i}$ are isomorphic for all $i = 1, 2, \dots, 5$. So D is isomorphic to one of

$$\{D_{i,j} | 1 \leq i \leq j \leq 5\}.$$

We can see these 15 digraphs are not isomorphic. \square

By Theorem 2, we can conclude that if D is a strongly connected 2-regular digraph and $\text{diam}(D) = d \geq 6$, then D has at least $d+3$ vertices and the extremal cases are given in Figure 2.

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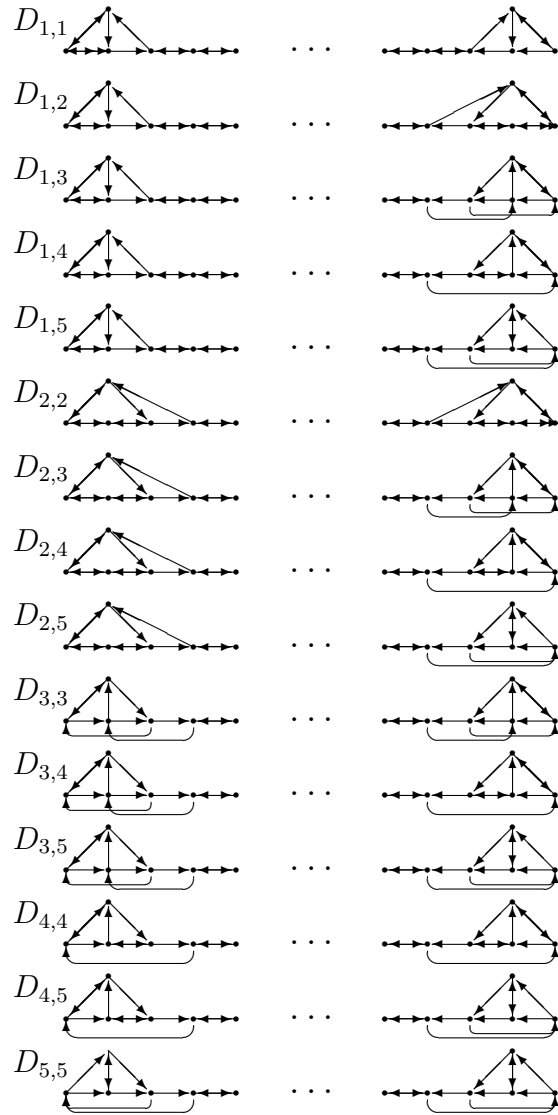


Figure 2. Extremal digraphs

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